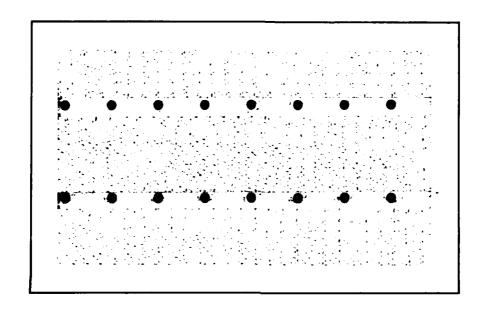


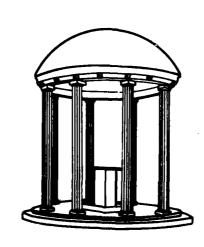
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VARIANCE REDUCTION IN THE SIMULATION OF STOCHASTIC ACTIVITY NETWORKS

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Abstract

This paper describes a Monte Carlo method based on the theory of quasirandom points for estimating the distribution functions and means of network completion time and shortest path time in a stochastic activity network. In particular, the method leads to estimators whose variances converge faster than 1/K, where K denotes the number of replications collected in the experiment. The paper demonstrates how accuracy diminishes for a given K with increasing dimensionality of the network and shows how a procedure that uses a cutset of the network together with convolution can reduce dimensionality and increase accuracy. Two examples illustrate the benefits of using quasirandom points together with a cutset and then convolution.

Key Words

Monte Carlo method

Networks

Quasirandom points

Variance reduction



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Introduction

Activity network analysis is a fairly standard tool in Operations
Research. When arc completion times in such a network are deterministic,
well-established algorithms exist for finding its network completion time and
shortest path time in a computationally efficient manner. For example, see
Elmaghraby (1977) and Wagner (1975). However, when arc completion times are
stochastic, analysis becomes considerably more difficult, even for relatively
small networks. Here the objective often is to compute the distribution
functions (d.fs.) and means of network completion time and shortest path
time, given the d.fs. of the statistically independent individual arc
completion times.

Because severe difficulties exist in deriving analytical solutions to these problems, many analysts have turned to Monte Carlo methods to derive approximate solutions. Van Slyke (1963) proposes the use of importance sampling to gain efficiency in estimating the characteristics of networks.

Martin (1965) describes how one can increase statistical efficiency when the arc passage time distributions are polynomials. Burt and Garman (1971) describe how conditional sampling can reduce dimensionality and thus sampling variation. Carrying this idea one step further Sigal, Pritsker and Solberg (1979, 1980) show how one can use a cutset of a network to effect similar reductions. The use of antithetic variates has also been proposed in a number of studies, including Cheng (1980), Sullivan, Hayya and Schaul (1982) and Grant (1982).

Although the benefits of these proposals are well established, they all lead to estimators with variances O(1/K) for K independent replications. The present paper shows how to increase the rate of convergence of this variance using quasirandom points and illustrates the success of the

approach for two networks of N=8 and N=18 arcs. The theory of quasirandom points has its origin in the numerical evaluation of multivariable integrals. The theory concerns the identification of vector sequences that when used in numerical integration lead to approximations whose absolute errors converge to zero faster than if independent random vectors were used. Since one can represent the desired d.fs. and means as multivariable integrals, the appeal of quasirandom points is apparent. Moreover, combining the use of quasirandom points with other approaches such as cutsets and convolution leads to even greater advantage, as we demonstrate.

Section 1 presents the stochastic network problem in detail. Section 2 describes how one can view the problem as one of multivariable numerical integration. Section 3 describes how one would tackle the problem by crude Monte Carlo methods and also describes the benefits of conditional sampling. Section 4 shows the benefits of using cutsets. Section 5 then indicates how cutsets together with convolution can significantly reduce the dimensionality of the problem.

Section 6 presents a short discussion of the principal concepts of multivariable numerical integration using quasirandom points and shows the extent to which known results apply to the problem at hand. Section 7 lays out an experiment whose design is used in Section 8 with two examples to show the benefit of using quasirandom points, especially when combined with convolution and the cutset approach.

Definitions

Consider an acyclic directed network with a single source, single sink, N arcs and L paths. Let X_1, \ldots, X_N , the passage times for arcs 1,...,N, be independent random variables where X_1 has distribution function (d.f.) F_1 on $[0,\infty)$ and inverse distribution function $G_1(u) = \min[x: F_1(x) > u$, 0 < u < 1]. The completion time of path m is

$$T_{m} = \sum_{i=1}^{N} a_{im} X_{i} = \sum_{i \in I_{m}} X_{i}$$
 (1)

where $a_{1m}=1$ if arc i is on path m , $a_{1m}=0$ otherwise and I_m denotes the set of arcs on path m .

The principal purpose of this study is to characterize the network completion time

$$T^* = \max(T_1, \dots, T_L) \tag{2}$$

and the shortest path time

$$T* = \min(T_1, \dots, T_L) . \tag{3}$$

For T*, characterizations include:

a.
$$pr(T^* < t)$$

0 < t < •

ъ. ЕТ*

pr(m is the longest path)

 $m = 1, \ldots, L$

d. pr(T* < t)

0 < t < •

e. E T*

f. pr(m is the shortest path)

n = 1,....L

The present study concerns the estimation of a, b, d and e . For expository convenience the main body of the paper concerns the estimation of $pr(T^* < t)$. The Appendix contains details for the estimation of the remaining quantities.

Let $|I_m|$ denote the cardinality of I_m . If

$$N = \sum_{m=1}^{L} |I_m| \tag{4}$$

then no paths share a common arc and T* and T* have d.fs.

$$F^{*}(t) = \prod_{m=1}^{L} F_{I_{m}}(t)$$
 (5)

and

$$F*(t) = 1 - \prod_{m=1}^{L} [1 - F_{I_m}(t)]$$
 $0 < t < \infty$

F1 being the convolution of the $|I_m|$ arcs on path m . As an example, m suppose that arc times are exponential with

$$F_i(t) = 1 - e^{-\lambda_i t}$$
 $\lambda_i > 0$ $i \in I_m$. (6)

Then one has for distinct λ_1

$$F_{\vec{l}_{m}}(t) = 1 - \sum_{\substack{i \in \vec{l}_{m} \\ j \neq i}} \frac{e^{-\lambda_{i}t}}{\prod_{\substack{j \in \vec{l}_{m}(1-\lambda_{i}/\lambda_{j})\\ j \neq i}}}.$$
(7)

Usually the summation in (4) exceeds N in which case the derivations of F^* and F_* are more complicated than in (5). In practice, even when (4) holds, the convolutions for F_{I_1}, \ldots, F_{I_L} may prove too complex to derive analytically.

2. A Solution via Integration

An alternative approach to characterization is through integration; in particular, numerical integration. Observe that one can write

$$F^{*}(t) = E I_{[0,t)}(T^{*})$$

$$= \int_{0}^{1} \cdots \int_{0}^{1} h(t;x_{1},...,x_{n}) \prod_{i=1}^{N} dF_{i}(x_{i})$$

$$= \int_{0}^{1} \cdots \int_{0}^{1} g(t;u_{1},...,u_{n}) du_{1},...,du_{N}$$
(8a)

where

$$I[a,b) (x) = 1$$
 if $a < x < b$
 $= 0$ otherwise,
 $h(t;x_1,...,x_N) = I[0,t) (max(\sum_{i=1}^{N} a_{in}x_i,...,\sum_{i=1}^{N} a_{iL}x_i))$ (9a)

and

$$g(t;u_1,...,u_N) = h(t;G_1(u_1),...,G_N(u_N))$$
 (9b)

Now (8b) involves multiple integration over the N-dimensional unit hypercube $[0,1)\times\ldots\times[0,1)$, hereafter denoted by V_N . It is the expression (9b) on which our analysis is based.

Consider the approximation

$$\bar{\mathbf{g}}_{K}(t) = \frac{1}{K} \sum_{j=1}^{K} \mathbf{g}(t; \mathbf{u}_{1j}, \dots, \mathbf{u}_{Nj})$$
 (10)

where $u_i = \{u_{i,j}; j = 1,2,...\}$ i=1,...,N are infinite sequences of points in V_N chosen in one of several ways. One way is through pure random sampling which we call the <u>crude Monte Carlo method</u>.

3. Crude Monte Carlo Sampling

Let

$$T_{j}^{*} = \max (T_{j1}, ..., T_{jL})$$
 (11)

where

$$T_{jm} = \sum_{i=1}^{N} a_{im} G_i(u_{ij})$$
 $m = 1,...,L$ (12)

and ujj i = 1,...,N j = 1,...,K are i.i.d. from U[0,1). Also define

$$g_j(t) = g(t; u_{1j}, \dots, u_{Nj}) = I_{[0,t)}(T_j^*)$$
 (13)

Then $g_{K}(t)$ in (10) has expectation $F^{*}(t)$ and

$$var \ \bar{g}_{K}(t) = \frac{F^{*}(t) [1 - F^{*}(t)]}{K} . \tag{14}$$

To improve on the result in (14), Burt and Garman (1971) introduce the use of conditional sampling. Suppose that there exist arcs i1,...,iL unique to paths 1,...,L respectively. Assume that these arcs are distinct and define

$$S_{jm} = \sum_{i=1}^{N} a_{im} G_i(u_{ij})$$

$$i \neq i_{m}$$
(15)

and

$$g_{j}^{(t)} = \prod_{m=1}^{L} F_{im}(t - S_{jm})$$
 (16)

Now $\mathbf{g}_{\mathbf{K}}(t)$ in (10), based on (16), has the correct expectation $\mathbf{F}^*(t)$, as can be seen by integrating (16) with respect to the joint p.d.f. of $\mathbf{S}_{\mathbf{j}_1},\ldots,\mathbf{S}_{\mathbf{j}_L}$. But $\mathbf{g}_{\mathbf{K}}(t)$ has smaller variance than (14), as a result of sampling from N-L instead of N arcs. Also note that

$$g_{j}(t) \leq \prod_{m=1}^{L} F_{1m}(t)$$
 (17)

provides an upper bound which also applies to $\bar{g}_{r}(t)$.

In principle, conditional sampling can be made more effective by working with subsets of arcs unique to each arc. Let

$$J_{m} = I_{m} \cap (\bigcup_{\substack{j=1 \ j \neq m}}^{L} I_{j})$$
 $m = 1, ..., L$ (18)

so that J_m is the set of arcs unique to path $\,m\,$. Assume $\,|\,J_m\,|\,>\,1\,$ m = 1,...,L and define

$$S_{jm} = \sum_{i \in I_m - J_m} a_{im} G_i(u_{ij})$$
 (19)

and

$$g_{j}(t) = \prod_{m=1}^{L} F_{j_{m}}(t-S_{jm})$$
 (20)

where F_J denotes the d.f. of the sum of the passage times on the arcs unique to path m. Since sampling in $N+L-\sum_{m=1}^{L}|J_m|$ dimensions is all that is necessary, the use of (20) in (10) preserves unbiasedness and reduces variance even further.

In essence, the effect of conditional sampling is to replace the coefficient of 1/K in (14) by one of smaller magnitude but in no way to change the convergence rate with regard to K as $K+\infty$. Also, note the upper bound

$$g_{j}(t) \leq \prod_{m=1}^{L} F_{J_{m}}(t)$$
 (21)

which is tighter than the bound in (17).

4. Cutsets

In practice it is conceivable that at least one path does not have a unique arc. That is, $|J_{\mathbf{m}}|=0$ for at least one path \mathbf{m} . Figure 1 illustrates such a case. Although conditional sampling as described by Burt

Insert Fig. 1 about here.

and Garman does not apply here, another more general approach proposed by Sigal, Pritsker and Solberg (1979, 1980) does. Let $\mathcal H$ denote a cutset of the network. A cutset is a set of arcs that connects a set of nodes $\mathcal W$ containing the source with a set of nodes $\mathcal W$ containing the sink. Also, assume that each path has only one arc in $\mathcal H$. If each arc in $\mathcal H$ points from $\mathcal W$ to $\mathcal W$, $\mathcal H$ is called a uniformly directed cutset (UDC).

Define

$$S_{jm} = \sum_{i=1}^{N} a_{im} G_i(u_{ij})$$
 (22)

and

$$Y_{ij} = \sup_{m=1,...,L} (a_{jm} S_{jm})$$
 ieff. (23)

Then define

$$g_{i}(t) = \Pi_{i \in H} F_{i}(t - Y_{i})$$
 (24)

Now $\vec{s}_{K}(t)$ in (10) based in (1) gives an unbiased estimate of $F^{*}(t)$

with smaller variance than crude Monte Carlo sampling since one is sampling in only N-|H| dimensions. Since a network may have more than one UDC, using the one with <u>maximal cardinality</u> gives the greatest reduction in dimensionality. However, the determination of this cutset for an arbitrarily large network is not trivial, being an NP complete problem. Nevertheless, among alternative known cutsets for a given network, the one with largest cardinality serves our purpose best. As an example of the use of cutsets note that (e_2, e_3, e_4, e_7) and (e_3, e_5, e_6, e_7) are UDCs with maximal cardinality 4 for the network of Fig. 1, whereas the cutset (e_1, e_2, e_3) with cardinality 3 is not.

As a second example, consider the network in Fig. 2, taken from Battersby (1970), which describes the steps encountered in the partial overhaul of a unit in an oil refinery. Table 1 shows the incidence matrix $\underline{a} = || a_{im} ||$

Insert Fig. 2 about here.

for arcs (rows) and paths (columns). Observe that arcs (e2, e3, e5, e6, e9, e10, e12) and arcs (e2, e3, e5, e9, e10, e12, e15) form two UDCs with maximal cardinality 7. However, the cutset (e2, e3, e5, e6 + e15, e9, e10, e12), formed by combining arcs e6 and e15 (see Fig. 2) has cardinality 8 and can be used to advantage to reduce dimensionality.

Insert Table 1 about here.

5. Further Dimensional Reduction

One can achieve at least one additional reduction in the dimensionality

of sampling by exploiting available knowledge about a network. Let

$$I_{\mathbf{m}} = I_{\mathbf{m}} - I_{\mathbf{m}} \cap H \tag{25}$$

where H denotes a UDC. Then define the subsets

$$J_{m_{1},...,m_{r}} = \bigcap_{j=1}^{r} I_{m_{j}} - \sum_{\substack{m=1 \ m\neq m_{1},...,m_{r}}}^{L} [(\bigcap_{j=1}^{r} I_{m_{j}}) \cap I_{m}]$$
(26)

Here J_{m_1, \dots, m_r} denotes the subset of arcs, not in H, that are uniquely on paths m_1, \dots, m_r . Suppose there are q such nonempty sets R_1, \dots, R_q .

Let F_{R_1} denote the d.f. of the sum of the $|R_1|$ arc passage times corresponding to the arcs in R_1 and let G_{R_1} denote the corresponding inverse d.f. Then define

$$S_{jm} = \sum_{i=1}^{q} (\Pi_{j \in R_i} a_{jm}) G_{R_i}(U_{ij})$$
 (27)

and

$$Y_{1j} = \sup_{m=1,...,L} (S_{jm})$$
 $i = 1,..., |H|$ (28)

and

$$g_{j}(t) = \prod_{i \in \mathcal{H}} F_{i}(t - T_{ij}) . \qquad (29)$$

Now, $\overline{g}_K(t)$ in (10) based on (29) again is an unbiased estimator of $F^*(t)$. However, since $N - |H| + q - \sum_{i=1}^q |R_i|$ dimensions arise for sampling per replication, one anticipates that var $\overline{g}_K(t)$ has smaller magnitude than previously described methods produce. Note that this reduced dimensionality comes from the use of UDCs together with convolution.

6. Multivariable Numerical Integration

As (8b) shows, one can view $F^*(t)$ as the result of a multivariable integration over a restricted region in V_N . We now describe how taking this

view of the problem leads to useful results when quasirandom points are used. Our account follows Kuipers and Niederreiter (1974) and Niederreiter (1978). Also see Schmidt (1973).

Suppose that our objective is to evaluate the integral

$$B = \int_{0}^{1} \dots \int_{0}^{1} f(x_{1}, \dots, x_{N}) dx_{1} \dots dx_{N}$$
 (30)

by

$$B_K = \frac{1}{K} \sum_{j=1}^{K} f(u_{1j}, \dots, u_{Nj})$$
 (31)

and to derive bounds on the error

$$\Delta_{K} = |B - B_{K}| . (32)$$

Three definitions facilitate our description. Consider the sequences

$$\{u_j = (u_{1j}, \dots, u_{Nj}); j = 1, \dots, K\}$$
.

<u>Definition 1.</u> A(R;K) = number of points u_1, \dots, u_K that fall in $R \subset V_N$.

Definition 2. u_1, \dots, u_K are uniformly distributed in V_N if

$$\lim_{K\to\infty}\frac{A(R;K)}{K\lambda(R)}=1$$

for all $R = \{(x_1, ..., x_N): \alpha_i < x_i < \beta_i \ i = 1, ..., N\}$ $R \subset V_N$ and where $\lambda(R)$ is the measure of volume of R in V_N .

Definition 3. The extreme discrepancy associated with a sequence u_1, \dots, u_K is

$$D_{K} = D_{K}(y_{1}, \dots, y_{K}) = \sup_{R^{*} \subset V_{N}} \left| \frac{A(R^{*}; K)}{K} - \lambda(R^{*}) \right|$$

where $K^* = \{(x_1, \dots, x_N): 0 \le x_1 \le \beta_1 \quad i = 1, \dots, N\}$. Hereafter we assume that $\{y_j \mid j = 1, \dots, K\}$ is a uniformly distributed sequence on V_N .

These definitions together with several fundamental theorems from the theory of quasirandom points provide us with a way of characterizing the error (32). In particular,

Theorem 1 (Hlawka 1961). If f is of bounded variation in the sense of Hardy and Krause, then

$$\Delta_{K} < V(f,N)D_{K} \tag{33}$$

where V(f,N) is a function of the bounded variation in f in N and lower dimensions.

Theorem 2 (van Aardenne-Ehrenfest 1945). For any infinite sequence u_1, u_2, \dots with N > 1

Theorem 3 (Roth 1954). For any sequence of K points in U_N with N > 2 $D_K > \frac{C_N \; (\log \; K)(N-1)/2}{K}$

where CN is a function of N only.

Theorem 4 (Roth 1954). For any infinite sequence in V_N with N > 1 $D_K > \frac{C_N (\log K)^{N/2}}{K}$

where Cy is a function of N only.

Theorem 1 shows that one can bound the error (32) by a quantity proportional to the extreme discrepancy D_K , a most convenient result. Theorem 2 implies that one cannot expect the extreme discrepancy to converge as fast as 1/K. Theorems 3 and 4 provide lower bounds on how fast convergence can occur.

We now describe a uniformly distributed sequence for which upper bounds are known. If R > 2 is an integer, then every non-negative integer n has an expansion of the form

$$n = \sum_{i=0}^{m} a_i R^i \qquad a_i \in \{0,1,\ldots,R-1\}$$

$$0 \le i \le m \text{ and } m = \lfloor \log_R n \rfloor.$$
(34)

Corresponding to this expansion one has the radical inverse function

$$\Phi_{R}(n) = \sum_{i=0}^{m} a_{i}R^{-i-1}$$
 (35)

We now define several sequences based on (34):

van der Corput Sequence: $\{\phi_2(j); j = 0,1,...,K-1\}$

Hammersley Sequence:
$$\{\phi_{R_1}(j), \ldots, \phi_{R_{N-1}}(j), \frac{j}{K}; j = 0,1,\ldots,K-1\}$$

where R1,..., RK-1 are pairwise relatively prime.

Halton Sequence:
$$\{\phi_{R_1}(j), \ldots, \phi_{R_N}(j); j=0,1,\ldots,K-1\}$$

where R1,...,RN are pairwise relatively prime.

The van der Corput sequence is uniformly distributed on V_1 and was the first sequence of this kind proposed for univariate numerical integration. The Hammersley and Halton sequences are uniformly distributed on V_N and as Theorems 5 and 6 show they offer useful upper bounds.

Theorem 5 (Halton 1960). For the Hammersley sequence

$$D_{K} < \frac{(\log K)^{N-1}}{K} \quad \prod_{i=1}^{N-1} \left(\frac{3R_{i}-2}{\log R_{i}} \right) \qquad N > \max_{i} R_{i}.$$

Theorem 6 (Halton 1960). For the Halton sequence

$$D_{K} < \frac{(\log K)N}{K}$$
 $\prod_{i=1}^{N-1} (\frac{3R_{i}-2}{\log R_{i}})$ $N > \max_{i} R_{i}$.

Several issues of significance need to be mentioned now. Firstly, since the van der Corput, Hammersley and Halton sequences are deterministic the upper bounds can be interpreted as worst case bounds with certainty. Secondly, although V(f,N) in (33) is in theory computable, in the present case f is unknown. Thirdly, the bounds presented here hold for integration the entire N-dimensional hypercube; however, as (8a), (9a) and (9b) show, our problem calls for integration over a restricted region in V_N . Now it is known (Niederreiter 1978, p. 982) that for integration over an arbitrary

region in VN

$$\Delta_{K} < C_{N}^{**}(D_{K})^{1/N} \tag{36}$$

where $C_N^{\prime\prime}$ is a function of N . This is a considerably weaker bound but since it applies to such a wider set of regions in V_N than (33) does there is reason to believe that the use of the Hammersley and Halton sequences can lead to accelerated convergence.

The fourth issue concerns boundedness. In the case of estimating $F^*(t)$ and $F_*(t)$, this assumption is met. However, it is not met for ET^* and ET_* . This fact together with the integration over a restricted region in V_N makes clear that when we move from theory to practice, some skepticism exists as to how efficient our approach will be. As the examples in Section 8 show, there is little basis for this skepticism.

The fifth issue relates to dimensionality. As Theorems 5 and 6 make clear, it is in one's best interest to reduce dimensionality as much as possible before using quasirandom points. The cutset approach together with convolution are the methods we use to effect this reduction.

The sixth issue of significance concerns the choice between the Hammersley (finite) and Halton (infinite) sequences in practice. If one uses the Hammersley sequence for fixed K and then decides that the accuracy is not acceptable, there is no recourse to continued exploitation of its special structure since $0,1/K,\ldots,(K-1)/K$ are the values in the Nth dimension. By contrast, the Halton sequence enables us to continue exploitation by merely generating additional points $\{\phi_{R_1}(j),\ldots,\phi_{R_N}(j)\}$ j=k, k+l,... The modest degradation in the upper bound that arises when using the Halton rather than

the Hammersley sequence seems worth the price.

7. Experimental Design

This section describes the layout for an experiment designed to determine the extent to which quasirandom points lead to accelerated accuracy when estimating $F^*(t)$, ET^* , $F_*(t)$ and ET_* for the networks in Figs. 1 and 2. In particular, we introduce a degree of randomness into the experiment in order to compute estimates of the variances of our point estimates. We then study the behavior of these sample variances as K increases.

resider an experiment consisting of Q statistically independent blocks or macroreplications each of K microreplications. Let $\{U_{\underline{i}\underline{m}}: i=1,...,N;$ $\underline{m}=1,...,Q\}$ denote a sequence of i.i.d. random variables from U[0,1) and define

$$S_{im} = (s: \phi_{R_i}(s) = U_{im})$$
 $i=1,...,N$ $m=1,...,Q$. (37)

Then on macroreplication m we use the quasirandom point

$$\phi_{R_1}(S_{1m} + j-1), \dots, \phi_{R_N}(S_{Nm} + j-1)$$
 on microreplication $j = 1, \dots, K$.

Let $\hat{\theta}_{jm}$ denote the unbiased estimate of a particular quantity θ computed on microreplication j on macroreplication m. Then $\bar{\theta}_1,\ldots,\bar{\theta}_Q$, where

$$\bar{\theta}_{m} = \frac{1}{K} \sum_{j=1}^{K} \hat{\theta}_{jm} \qquad m=1,\dots,Q, \qquad (38)$$

are i.i.d. random variables with sample variance

$$a_{K}^{2} = \frac{1}{Q-1} \sum_{m=1}^{Q} (\vec{\theta}_{m} - \vec{\theta}_{Q})^{2}$$
 (39)

where

$$\bar{\theta}_0 = \frac{1}{0} \sum_{m=1}^{Q} \bar{\theta}_m$$
 .

It is the behavior of so versus K that interests us.

To provide an instructive normalization we also run M statistically independent macroreplications each of K=1 microreplication. Let $\hat{\theta}_1, \ldots, \hat{\theta}_M$ denote the resulting estimates of θ each with sample variance

$$w_{Q}^{2} = \frac{1}{M-1} \sum_{m=1}^{M} (\hat{\theta}_{m} - \theta_{Q})^{2}$$
 (40)

where

$$\tilde{\theta}_{m} = \frac{1}{M} \sum_{m=1}^{M} \hat{\theta}_{m}$$
.

Then the quantity w_M^2 /Ks $_K^2$ should increase as K increases if accelerated convergence is occurring.

We set $M=10^5$. For the quasirandom case we perform runs for $K=2^{\frac{1}{3}}$, $j=1,2,\ldots$ and set $Q=2^{16}/K$ for $K<2^9$ and $Q=2^7$ for $K>2^9$. Following this procedure enables us to treat s_K^2 and w_M^2 as highly reliable point estimators of the quantities var $\overline{\theta}_m$ and var $\widehat{\theta}_m$ respectively.

8. Examples

We refer to the network in Fig. 1 as the SPS network and to the one in Fig. 2 as the Battersby network. Table 2 lists the relevant cutsets and noncutsets. In both examples the arc passage times are assumed independent

Insert Table 2 about here.

and exponentially distributed. For the SPS network the means are $\lambda_{1}=10$, $\lambda_{2}=15$, $\lambda_{3}=18$, $\lambda_{4}=6$, $\lambda_{5}=7$, $\lambda_{6}=3$, $\lambda_{7}=10$ and $\lambda_{8}=2$. For arc i on microreplication j on macroreplication m the arc passage time is

$$X_i = G_i(\phi_{R_i}(S_{im} + j-1)) = -\lambda_i \ln(1 - \phi_{R_i}(S_{im} + j-1))$$
 (41)

For the runs using quasirandom points, the cutset $H=(e_2, e_3, e_4, e_7)$ was used reducing dimensionality to N-|H|=8-4-4. An algorithm due to Halton

and Smith (1964) was used to compute the Halton sequences. For the M independent microreplications all N arc times from (41) were used.

Table 3 shows the behavior of w_M^2/Ks_K^2 for $F^*(t)$ estimated at t=15, 30, 40, 50, 70 for ET^* , for $F_*(t)$ estimated at t=5, 10, 15, 20, 25 and for ET_* . As a point for comparison, Sigal, Pritsker and Solberg (1979)

Insert Table 3 about here.

report variance reductions ranging from 4.18 to 11.00 for the estimation of $F^*(t)$ at t=15, 30, 40, 50 and 70 using the cutset μ and independent microreplications.

The second example studies variance reduction for the Battersby network using the means in Table 1. Two separate sets of experiments were run. The first set uses partition 1 in Table 2 with the cutset $\mu=(e_2, e_3, e_5, e_6, e_9, e_{10}, e_{12})$ with quasirandom points, yielding a dimensionality of $N-|\mu|=11$. The second set uses partition 2 with the cutset $\mu=(e_2, e_3, e_5, e_6+e_{15}, e_9, e_{10}, e_{12})$ yielding a dimensionality of 6 after appropriate convolution within the noncutset. The quantity $F^*(t)$ was estimated for t=90, 110, 120, 130, 140, 150, 160, 180, 190 and $F^*(t)$ for t=20, 25, 30, 35, 40, 50, 60, 80, 90. Table 4 shows the resulting variance reductions V_R^2/Ks_R^2 for $F^*(t)$, ET^* , $F^*(t)$ and ET^* .

Insert Table 4 about here.

Tables 3 and 4 admit several instructive observations:

- As K increases variance reduction increases; however,
 variance reduction is not necessarily monotone in K.
- 4. For given K, variance reduction tends to decline with increasing dimensionality.
- 3. For given K variance reduction is greater for F*(t) than for ET* and for F*(t) than for ET*, showing the benefit one obtains from bounded integrands.

Although these observations establish the benefits of using quasirandom points with the cutset approach and convolution, there is a cost issue that also needs to be considered. If one compares the cost of producing one sample point for $F^*(t)$, ET^* , $F_*(t)$ and ET_* using quasirandom points to the cost of producing a sample point using pure random sampling, the ratio never exceeds 2.5 in our experiments. The most costly step turns out to be sampling the summed arc times $x_1 + x_2$, in the noncutset of partition 2 in Table 2, from the distribution function

$$F(t) = 1 - \frac{e^{-\lambda_1 t}}{1 - \lambda_1 / \lambda_1} - \frac{e^{-\lambda_1 t}}{1 - \lambda_1 / \lambda_1}$$

This was done using the Newton-Raphson method.

Since this cost ratio is independent of K, we have strong evidence that for sufficiently large K, the quasirandom point method offers a clear advantage over pure random sampling for networks of arbitrary size. A program called NETWOK is currently under development to implement the proposed method in a computationally efficient manner.

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Appendix

Estimation of F*(t). We derive the estimator based on the cutset H and (22). Let

$$Z_{ij} = \inf_{m=1,\dots,L} (S_{jm})$$
 ieff $a_{im}=1$

and define

$$h_{j}(t) = 1 - \pi_{i \in H} [1 - F_{i}(t - Z_{ij})]$$
.

Then

$$\vec{h}_K(t) = \frac{1}{K} \sum_{j=1}^{K} h_j(t)$$

gives an unbiased estimator of F*(t) based on sampling from N - |H| dimensions.

Estimation of ET* and ET*. Using (11) one has the unbiased estimator

$$\vec{T}_{K}^{\star} = \frac{1}{K} \sum_{j=1}^{K} T_{j}^{\star}$$

for ET* . Let

$$T*_j = min(T_{j1}, \dots, T_{jL})$$
.

Then

$$\bar{\mathbf{T}}_{K} = \frac{1}{K} \sum_{j=1}^{K} \mathbf{T}_{j}$$

is an unbiased estimator of ET* .

TABLE 1
Incidence Matrix for Battersby Network

| | | l | | | | Pa | th m | | | |
|-----------------|----|---|---|-----|-----|----|------|-----|---|---|
| Arc i | λį | | 1 | . 2 | . 3 | 4 | 5. | . 6 | 7 | 8 |
| el | 16 | | 1 | 1 | | | | | | |
| e ₂ | 16 | | 1 | | | | | | | |
| ез | 8 | | | 1 | | | | | | |
| e 4 | 6 | | 1 | 1 | | | | | | |
| e 5 | 16 | | | | 1 | 1 | | | | |
| e 6 | 40 | | | | | | 1 | | | |
| e 7 | 24 | | | | | | | 1 | 1 | 1 |
| eg | 16 | | | | | | | | 1 | 1 |
| e 9 | 16 | 1 | | | | | | | | 1 |
| e 10 | 24 | | | | | | | | 1 | |
| e 11 | 8 | | | | | | | | 1 | 1 |
| e ₁₂ | 4 | | | | | | | 1 | | |
| e13 | 36 | | 1 | 1 | 1 | | | | | |
| e 14 | 12 | | | | | 1 | | | | |
| e ₁₅ | 8 | | | | | | 1 | | | |
| e16 | 24 | | | | | | | 1 | 1 | 1 |
| e 17 | 8 | | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| e 18 | 24 | | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |

Table 2
Network Analysis

| Network | Cutset | Noncutset | | |
|-----------------------------------|---------------------------------------|---|--|--|
| Fig. 1 | ^c 2, e3, e4, e7 | e ₁ , e ₅ , e ₆ , e ₈ | | |
| Battersby (Fig. 2) Partition 1 | e2, e3, e5, e6, e9, e10, e12 | e1, e4, e7, e8, e11, e13, e14, e15, e16, e17, e18 | | |
| Battersby (Fig. 2) Partition 2 | e2, e3, e5, e6 + e15, e9, e10, e12 | e1 + e4, e7 + e16, e8 + e11, e13, e14, e17 + e18 | | |

 $e_i + e_j \equiv convolution of the d.fs. of arcs i and j.$

M=105 , N=8 , L=6 , |H|=4

| K | F* | (t) | ET* | F: | *(t) | ET* |
|-----|-------|--------|-------|--------|--------|-----|
| | min | max | | min | max | |
| 2 | 42 | 106 | 7 | 46 | 196 | 8 |
| 22 | 53 | 96 | 9 | 44 | 345 | 9 |
| 23 | 74 | 101 | 12 | 53 | 773 | 11 |
| 24 | 87 | 177 | 17 | 125 | 957 | 19 |
| 25 | 111 | 258 | 23 | 149 | 1,285 | 24 |
| 26 | 146 | 701 | 27 | 313 | 2,722 | 36 |
| 27 | 248 | 1,066 | 51 | 534 | 2,984 | 45 |
| 28 | 365 | 1,409 | 83 | 710 | 3,183 | 51 |
| 29 | 383 | 2,426 | 130 | 1,527 | 5,323 | 123 |
| 210 | 987 | 4,612 | 153 | 2,197 | 4,512 | 138 |
| 211 | 1,000 | 5,221 | 266 | 3,438 | 8,604 | 260 |
| 212 | 1,518 | 11,341 | 338 | 5,391 | 13,213 | 303 |
| 213 | 1,978 | 16,907 | 548 | 17,299 | 39,357 | 484 |
| 214 | 3,289 | 51,709 | 1,258 | 27,943 | 63,234 | 838 |

[†]min and max for $F^*(t)$ denote, respectively, the minimal and maximal variance reduction for t=15, 30, 40, 50, 70 · min and max for $F_*(t)$ denote, respectively, the minimal and maximal variance reduction for t=5, 10, 15, 20, 25 ·

Table 4

Variance Reduction $\frac{w^2}{M}/\frac{g^2}{g}$ for the Battersby Network†

| | | Partition | tion 1 | | | M=10 ⁵ , N= | N=18 , L=8 | | Partition 2 | 2 | | |
|-----|------------|-----------|----------|-----|----------|------------------------|------------|-------|-------------|-------|-------|-----|
| ۵. | # N 1 | P*(t) | ET.* | E : | F*(t) | ET* | ** | F*(t) | ET* | ₹. | F*(t) | ET. |
| | | Y S | | | ERX | | 11 | Tagx | | ala | E8X | |
| 7 | ∞ | 11 | . | 5 | ∞ | 4 | 12 | 18 | 9 | 9 | 11 | 2 |
| 22 | . | 7 | m | က | 4 | 2 | 13 | 18 | 9 | \$ | 6 | 4 |
| 23 | 4 | Ŋ | m | 7 | ю | 2 | 19 | 25 | 7 | 7 | 14 | 9 |
| 24 | 'n | m | 7 | - | 7 | 1 | 24 | 46 | 6 | 22 | 26 | 11 |
| 25 | 1 0 | 7 | 4 | 1 | m | 7 | 28 | 54 | 16 | 22 | 76 | 12 |
| 56 | 11 | 23 | 11 | 'n | 11 | 18 | 38 | 92 | 16 | 56 | 86 | 18 |
| 27 | 14 | 32 | 15 | က | 14 | 16 | 62 | 153 | 22 | 65 | 156 | 36 |
| 28 | 13 | 38 | 14 | 7 | 7 | 10 | 81 | 226 | 19 | 81 | 167 | 44 |
| 29 | 15 | 38 | 16 | 7 | 7 | 11 | 100 | 332 | 31 | 223 | 330 | 88 |
| 210 | 17 | 38 | 15 | 7 | 11 | 10 | 203 | 809 | 26 | 297 | 619 | 110 |
| 211 | 72 | 109 | 27 | 19 | 33 | 16 | 473 | 1,164 | 104 | 867 | 1,178 | 395 |
| 212 | 29 | 176 | 84 | 10 | 59 | 112 | 637 | 1,864 | 186 | 1,064 | 2,743 | 411 |
| 213 | 115 | 250 | 11 | 10 | 78 | 06 | | | | | | |

Thin and hax for P*(t) denote, respectively, the minimal and waximal variance reduction for t=90,110,120,130,140,150,160,180,190. Hin and hax for F*(t) denote, respectively, the minimal and maximal variance reduction for t=20,25,309,35,40,50,60,80,90.

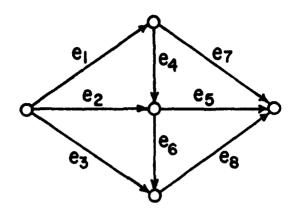


Figure 1. SPS Network

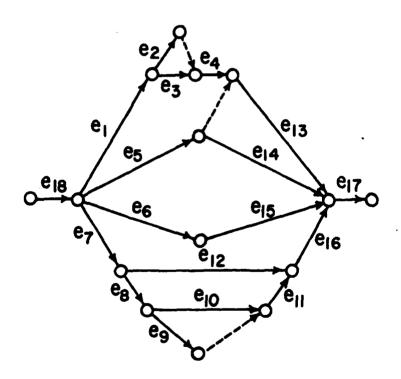


Figure 2. Battersby Network

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| This paper describes a Monte Carlo method based on the theory of | | | | | | |
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| network. In particular, the method leads to estimators whose variances | | | | | | |

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converge faster than 1/K, where K denotes the number of replications collected in the experiment. The paper demonstrates how accuracy diminishes for a given K with increasing dimensionality of the network and shows how a procedure that uses a <u>cutset</u> of the network together with convolution can reduce dimensionality and increase accuracy. Two examples illustrate the benefits of using quasirandom points together with a cutset and then convolution.

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